

# A Timer-Based Distributed Channel Access Mechanism in Networked Control Systems

**Tahmoores Farjam<sup>\*</sup>**, T. Charalambous<sup>\*</sup> and H. Wymeersch<sup>#</sup>

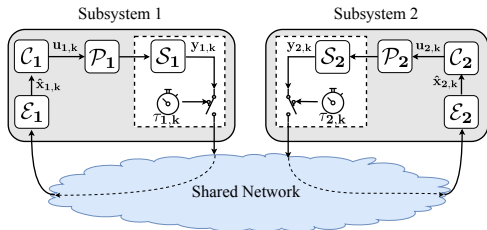
<sup>\*</sup> School of Electrical Engineering, Aalto University

<sup>#</sup>Department of Electrical Engineering, Chalmers University of Technology

ISCAS 2018  
Florence, Italy

## Networked Control Systems

- The NCS consists of  $N$  dynamical subsystems sharing the limited communication resources
- Each subsystem contains **smart sensors** which are capable of local computations
- Local measurements are sent through the network to be received by their corresponding estimator



## System and Network Models

- Each subsystem is modeled by a discrete-time LTI stochastic process

$$\begin{aligned}x_{i,k+1} &= A_i x_{i,k} + B_i u_{i,k} + w_{i,k}, \\ y_{i,k} &= C_i x_{i,k} + v_{i,k},\end{aligned}$$

- Variable  $\delta_{i,k}$  denotes whether the measurements of subsystem  $i$  is transmitted or not

$$\delta_{i,k} = \begin{cases} 1, & y_{i,k} \text{ is transmitted,} \\ 0, & \text{otherwise.} \end{cases}$$

- The communication channels are assumed to be reliable, therefore if no collision happens

$$\delta_{i,k} = 1 \rightarrow y_{i,k} \text{ is received at destination}$$

## Estimator

- Define

$$\hat{x}_{i,k+1|k} = \mathbb{E}\{x_{i,k+1}|\mathcal{Y}_k\},$$

$$P_{i,k+1|k} = \mathbb{E}\{(x_{i,k} - \hat{x}_{i,k+1|k})(x_{i,k} - \hat{x}_{i,k+1|k})^T|\mathcal{Y}_k\},$$

$$\hat{x}_{i,k+1|k+1} = \mathbb{E}\{x_{i,k+1}|\mathcal{Y}_{k+1}\},$$

$$P_{i,k+1|k+1} = \mathbb{E}\{(x_{i,k+1} - \hat{x}_{i,k+1|k+1})(x_{i,k+1} - \hat{x}_{i,k+1|k+1})^T|\mathcal{Y}_{k+1}\}.$$

- Kalman filter [Sinopoli *et al.*, 2004]

$$\hat{x}_{i,k+1|k} = A_i \hat{x}_{i,k|k} + B_i u_{i,k},$$

$$P_{i,k+1|k} = A_i P_{i,k|k} A_i^T + W_i,$$

$$K_{i,k+1} = P_{i,k+1|k} C_i^T (C_i P_{i,k+1|k} C_i^T + V_i)^{-1},$$

$$\hat{x}_{i,k+1|k+1} = \hat{x}_{i,k+1|k} + \delta_{i,k+1} K_{i,k+1} (y_{i,k+1} - C_i \hat{x}_{i,k+1|k}),$$

$$P_{i,k+1|k+1} = (I - \delta_{i,k+1} K_{i,k+1} C_i) P_{i,k+1|k}.$$

## Controller

- The cost function is defined as

$$J_{i,0} = \mathbb{E} \left\{ \sum_{k=0}^{\infty} (x_{i,k}^T Q_i x_{i,k} + u_{i,k}^T R_i u_{i,k}) \right\}$$

- The control law which minimizes this cost is

$$u_{i,k} = L_i \hat{x}_{i,k|k}$$

where

$$L_i = -(B_i^T \Pi_i B_i + R_i)^{-1} B_i^T \Pi_i A_i$$

$$\Pi_i = A_i^T \Pi_i A_i - A_i^T \Pi_i B_i (B_i^T \Pi_i B_i + R_i)^{-1} B_i^T \Pi_i A_i + Q_i$$

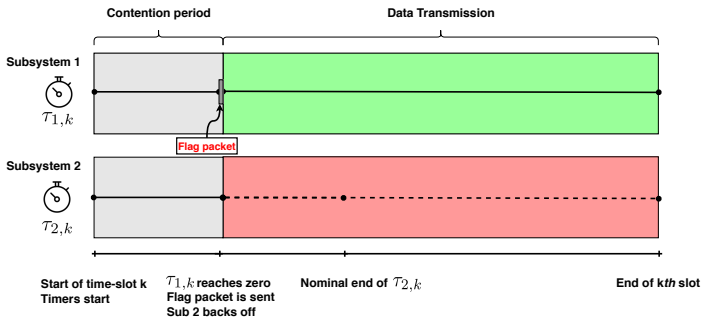
- The feedback gain  $L_i$  does not depend on scheduling or characteristics of the disturbances

## Concept of Local Timers

- Every subsystem possesses a local timer which is set to

$$\tau_{i,k} = \frac{\lambda}{m_{i,k}}$$

- The timers start at the beginning of each time slot and the subsystem whose timer finishes first begins transmission after sending a flag packet

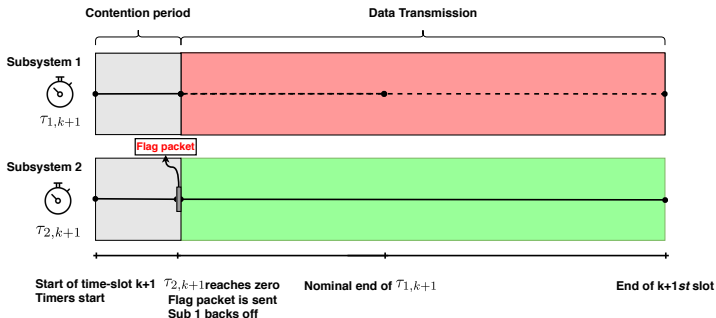


## Concept of Local Timers

- Every subsystem possesses a local timer which is set to

$$\tau_{i,k} = \frac{\lambda}{m_{i,k}}$$

- The timers start at the beginning of each time slot and the subsystem whose timer finishes first begins transmission after sending a flag packet



## Cost of Information Loss (CoIL)

- Minimizing the quadratic cost at each step  $k$  is equivalent to

$$\text{minimize } \sum_{i=1}^N \text{tr} (\Gamma_{i,k} P_{i,k|k}),$$

where  $\Gamma_i = L_i^T (B_i^T \Pi_i B_i + R_i) L_i$ .

- According to Kalman filter equations

$$P_{i,k|k} = \begin{cases} (I - K_{i,k} C_i) P_{i,k|k-1}, & \delta_k = 1 \\ P_{i,k|k-1}, & \delta_k = 0 \end{cases}$$

Hence, the increase in total cost in case subsystem  $i$  does not transmit at step  $k$  is defined as CoIL [Charalambous *et al.*, 2017]

$$\text{CoIL}_{i,k} = \text{tr} (\Gamma_i (P_{i,k|k-1} - P_{i,k|k})),$$



## Timer-based Mechanism (TBCoIL)

---

### Algorithm 1: TBCoIL for general case of $c$ available channels

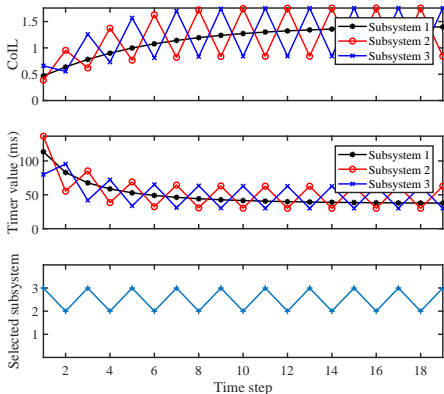
---

**Input:**  $c, \lambda$

- 1 Initialization: Start timers from  $\tau_j = \frac{\lambda}{C_{\text{CoIL}}}$  and let number of flags  
 $f = 0$
- 2 **while**  $f < c$  **do**
- 3     **for**  $j \leftarrow 1$  **to**  $c$  **do**
- 4         **if**  $\tau_j \neq 0$  *and timer is running* **then**
- 5             listen for flags
- 6             **if** *flag is received* **then**
- 7                 freeze  $\tau_j$  and  $f = f + 1$
- 8             **end**
- 9         **else if**  $\tau_j = 0$  **then**
- 10             send flag and set  $f = c$
- 11             freeze all running timers
- 12         **end**
- 13     **end**
- 14 **end**
- 15 **if** *any*  $\tau_j = 0$  **then** transmit on channel  $j$

## Illustrative Example

- 3 subsystems have access to 1 channel
- subsystem 1 is stable while 2 & 3 are unstable
- $\lambda = 52.82 \times 10^{-3} m^2 s$

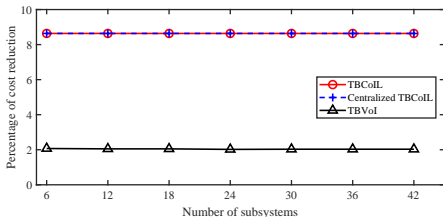


## Numerical Results 1

- The subsystems are selected from two homogeneous classes of unstable (I) and stable (II)

$$A_I = \begin{bmatrix} 1.1 & 0 \\ 0 & 1 \end{bmatrix}, A_{II} = \begin{bmatrix} 0.9 & 0 \\ 0 & 0.9 \end{bmatrix}, B = C = I_{2 \times 2}.$$

- Only 1/2 subsystems can transmit their measurements at each step
- Round-robin protocol is used as the basis of comparison

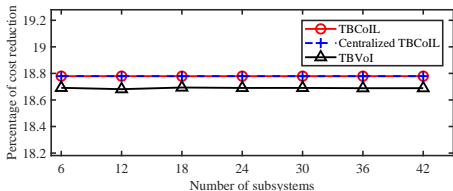


## Numerical Results 2

- The subsystems are selected from two homogeneous classes of stable (I) and unstable (II)

$$A_I = \begin{bmatrix} 1.4 & 0 \\ 0 & 1.2 \end{bmatrix}, A_{II} = \begin{bmatrix} 0.9 & 0 \\ 0 & 0.9 \end{bmatrix}, B = C = I_{2 \times 2}.$$

- Only 1/2 subsystems can transmit their measurements at each step
- Round-robin protocol is used as the basis of comparison



## Summary

---

- Novel distributed channel access mechanism was introduced
  - Concept of local timers for general case of prioritization
- CoIL was used as the local measure for implementation
- Distributed implementation was shown to perform as well as the optimal central scheduler

## Future directions

- Possibility of collision due to non-negligible duration of flag packets
- Extension to the case of wireless communication and presence of imperfect channels

Thank You!

---



**Questions?**